

Comparing Transparent Static Analyzers with Open Verification Dashboard

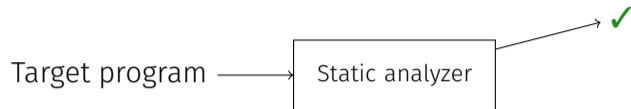
Tom Goalard, Karoliine Holter, Simmo Saan,
Vesal Vojdani, Raphaël Monat

`rmonat.fr/ecoop26_tsa`

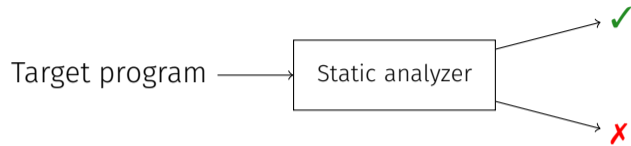
Target program



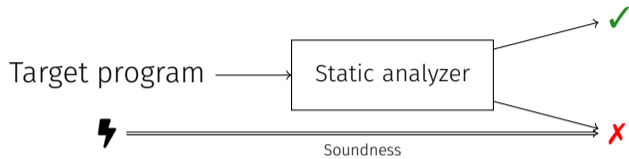
Static Analysis

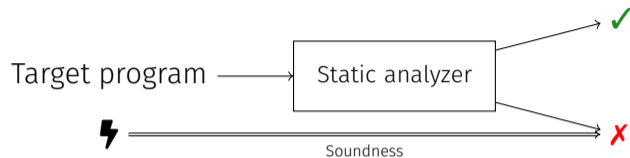


Static Analysis



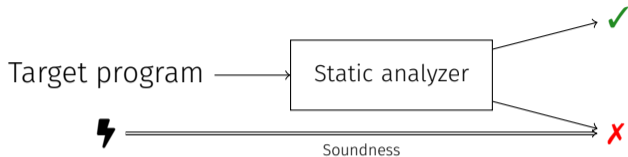
Static Analysis





Alarms

- ▶ Muske and Serebrenik [MS23]: 130 works on alarm postprocessing
- ▶ Alarms are only one side of the coin!



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Transparency

- ▶ Log proof obligations (PO). Analogy with deductive verification.
- ▶ Enables semantics-directed, fine-grained output.

Comparing Static Analyzers

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- `strcpy` misuse: buffer overflow or contract violation for the stub?

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 - ▶ Incomparable approximations between analyzers
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Dashboard Home Projects

```
1 int main_i;
2 void main() {
3   int a[0];
4   while (a[main_i])
5     main_i = main_i + 1;
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[Back to file list](#) Filter errors Error Types

reduced/standard_find_ground-1.i

Invalid memory access: 4.9-4.18
Error Level: The two analyzers disagree on the error level of this conflict

GOBLINT	MOPSA
Warning 4.9-4.18 Invalid array access: May access out of bounds	Error 4.9-4.18 accessing 4 bytes at offset 0 of variable 'a:./dashboard/examples/reduced/standard_find_ground-1.i:3.2-10' of size 0 bytes

Integer overflow: 4.9-4.18
Only One Proof Obligation: Only the first analyser has a proof obligation for this conflict

GOBLINT	MOPSA
Safe 4.9-4.18 Cast: true	No checks for this range

- 1 Transparent Static Analysis
- 2 Open Verification Dashboard
- 3 Experimental Evaluation

Transparent Static Analysis

Safety check categories

Example on toy imperative language with mathematical integers

$$\mathcal{C} = \{\text{divByZero}\}$$

\mathcal{L}

$$\varphi \in \text{PO} = \mathcal{L} \times \mathcal{C}$$

$$\Theta \subseteq \text{CH} = \text{PO} \times \{\checkmark, \times\}$$

Safety check categories

Program locations

Proof obligations

Checks

$$\mathbb{E}[\![e \in \text{Expr}]\!] : (\text{Var} \rightarrow \mathbb{Z}) \rightarrow 2^{\mathbb{Z}}$$

Concrete semantics of e

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Concrete semantics of e

$$\mathbb{E}[[e_1 /^\ell e_2]]\sigma =$$

$$\text{let } M_1 = \mathbb{E}[[e_1]]\sigma \text{ and } M_2 = \mathbb{E}[[e_2]]\sigma \text{ in}$$

$$\{m_1 / m_2 \mid m_i \in M_i, m_2 \neq 0, m_1 / m_2 \in \mathbb{Z}\}$$

$$\mathbb{E}[[e \in \text{Expr}]] : (\text{Var} \rightarrow \mathbb{Z}) \rightarrow 2^{\mathbb{Z}} \times 2^{\text{CH}}$$

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let $\Theta' = \{(\ell, \text{divByZero}), \text{check}(m_2 \neq 0) \mid m_2 \in M_2\}$ in

$\{m_1 / m_2 \mid m_i \in M_i, m_2 \neq 0, m_1 / m_2 \in \mathbb{Z}\}, \Theta_1 \cup \Theta_2 \cup \Theta'$

$\text{check}(b) = \text{if } b \text{ then } \checkmark \text{ else } \times$

$\mathbb{E}[[e \in \text{Expr}]] : (\text{Var} \rightarrow \mathbb{Z}) \rightarrow 2^{\mathbb{Z}} \times 2^{\text{CH}}$ *Concrete semantics of e*

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$\mathbb{E}[[1 /^\ell \text{rand}(-1, 1)]]\emptyset = \{-1, 1\}, \{(l, \text{divByZero}), \checkmark\} \cup \{(l, \text{divByZero}), \times\}$

$$\mathbb{E}^\sharp[[e \in \text{Expr}]] : (\text{Var} \rightarrow \mathbb{Z}^\sharp) \rightarrow \mathbb{Z}^\sharp$$

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$\Theta \subseteq \Theta^\#$ concisely encodes:

- ▶ If a check fails at runtime, the analysis should say so

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$\Theta \subseteq \Theta^\#$ concisely encodes:

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- ▶ If a check is reachable at runtime, the analysis should check it *safe check*, *false alarm*

From analysis checks (same location, multiple contexts) to analysis results:

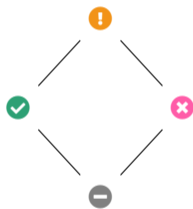
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$\emptyset \mapsto \ominus$

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Report postprocessing

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Open Verification Dashboard

Cross-Analyzer Comparison

Status ₁	Status ₂			
	!	×	✓	-
!	A ⁻	P	P	D
×	P	A ⁻	C	D
✓	P	C	A ⁺	D
-	D	D	D	A ⁺

► Positive or negative **agreement** (A⁺ / A⁻)

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- ▶ **Contradiction** (C)

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 - Detailed view

Open Verification Dashboard GUI – Detailed view

Dashboard

Home Projects

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[Back to file list](#) Filter errors Error Types

reduced/standard_find_ground-1.i

Invalid memory access: [4.9-4.18](#)

Error Level: The two analyzers disagree on the error level of this conflict

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Integer overflow: [4.9-4.18](#)

Only One Proof Obligation: Only the first analyser has a proof obligation for this conflict

GOBLINT	MOPSA
Safe 4.9-4.18 Cast: true	No checks for this range

Integer overflow: [5.4-5.23](#)

Safety W2: Only the second analyser says that this is safe

GOBLINT	MOPSA
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Proof Obligation Outcomes

- ▶ Report successful checks in addition to alarms
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Following a systematic comparison, 6 categories identified:

(signed|unsigned) integer overflow in (arithmetic operator|explicit cast|implicit cast)

Integrating a New Analyzer

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Standardized Integer Overflow Categories

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Program Range Reporting

Multiple choices (line only, line+column, line+columns) in IKOS, Goblint, Mopsa

↪ Heuristic location alignment (details in paper)

Experimental Evaluation

Benchmarks

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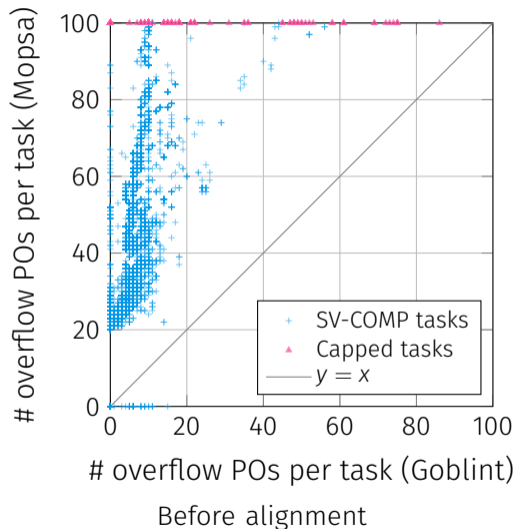
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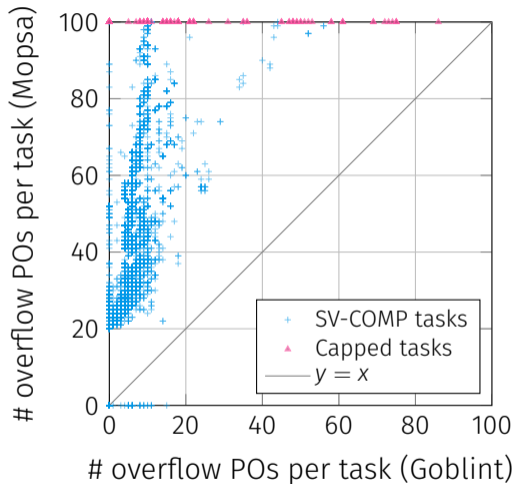
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- 3 Distillation of Disagreements through Testcase Reduction?

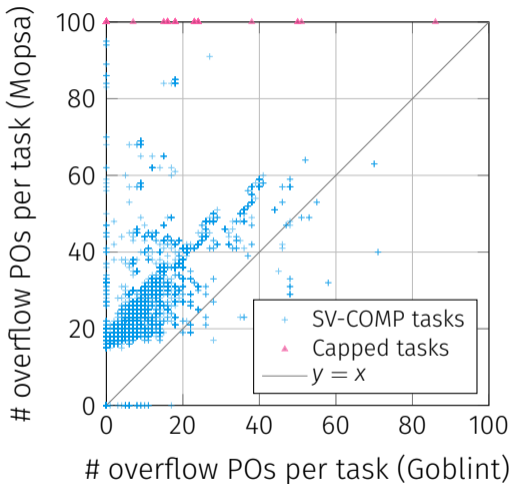
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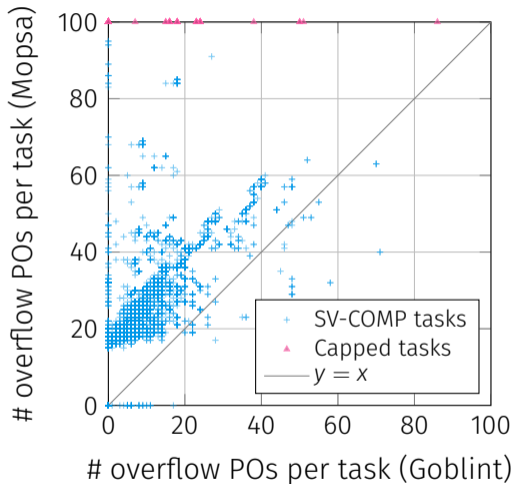
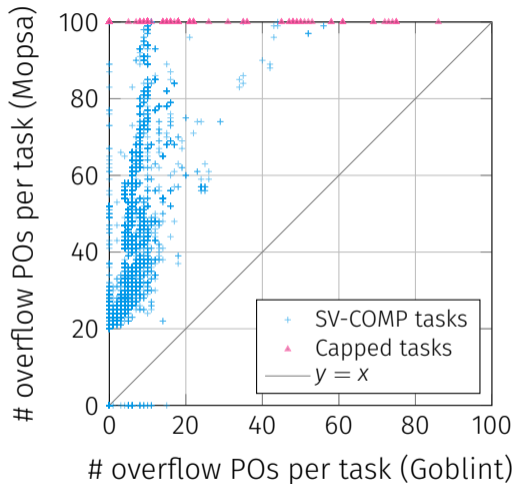


Before alignment



After alignment

SV-COMP Benchmark Validity?



N.B: Mopsa performs callstack-sensitive PO reporting.

RQ1: Complementarity

Mopsa	Goblint			
	!	×	✓	−
!	8122	0	2268	159
×	1	0	0	6
✓	222	0	11700	23754
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Issue	Mopsa	Goblint
Missing check		#1933, cil#215, cil#216
Wrong location	#248	cil#211
Spurious check	#251	#1909, #1910, #1932

RQ2: Joint Benefit (Coreutils uniq)

		Goblint			
Mopsa	!	×	✓	—	
!	16	0	3	0	
×	0	0	0	0	
✓	8	0	46	935	
—	2	0	14	0	

RQ2: Joint Benefit (Coreutils uniq)

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Joint Benefit

- ▶ Virtual combination improves precision

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Joint Benefit

- ▶ Virtual combination improves precision
- ▶ But currently insufficient to prove whole program correct

Automated testcase reduction oracles

- ▶ Verdict-based: boolean output of SV-COMP verification

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- ▶ Dashboard-based: specific safety check category

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Folder	Task Metrics			Verdict-Based			Dashboard-Based	
	#	Lines	Conflicts	Final lines	Pres. (%)	Time (s)	Final lines	Time (s)
float-benchs	6	196.17	1.00	8.67	100.00	184.50	8.67	233.17
nla-digbench	4	42.25	5.75	5.00	0.00	104.25	6.00	123.00
termination-crafted	11	26.18	1.55	3.73	18.18	119.73	5.91	156.45
termination-crafted-lit	11	27.18	1.73	5.27	45.45	92.00	7.00	102.73
termination-numeric	3	30.00	1.00	4.00	0.00	124.00	5.67	190.00
termination-restricted-15	13	21.62	1.00	4.31	7.69	83.62	8.77	98.62

Automated testcase reduction oracles

- ▶ Verdict-based: boolean output of SV-COMP verification
- ▶ Dashboard-based: specific safety check category

Folder	Task Metrics			Verdict-Based			Dashboard-Based	
	#	Lines	Conflicts	Final lines	Pres. (%)	Time (s)	Final lines	Time (s)
float-benchs	6	196.17	1.00	8.67	100.00	184.50	8.67	233.17
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Dashboard Strategy

- ▶ Less aggressive in reducing size,
- ▶ 100% preservation of disagreements.

Conclusion

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- ▶ Exchanging results: SARIF [OAS20], SV-COMP's witnesses [Bey+16; Aya+24]

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